

Key Ideas in Linear Algebra

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Multiply $A\mathbf{x}$ by columns, not rows

$$A\mathbf{x} = \begin{bmatrix} | & & | \\ \mathbf{a}_1 & \cdots & \mathbf{a}_n \\ | & & | \end{bmatrix} \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix} = \begin{bmatrix} \mathbf{a}_1 \end{bmatrix} x_1 + \cdots + \begin{bmatrix} \mathbf{a}_n \end{bmatrix} x_n$$

$A\mathbf{x}$ is a combination of the columns \mathbf{a}_1 to \mathbf{a}_n

Column space $\mathbf{C}(A) = \mathbf{all}$ combinations of the columns

$A\mathbf{x} = \mathbf{b}$ has at least one solution \mathbf{x} when \mathbf{b} is in $\mathbf{C}(A)$

Matrix multiplication: Columns times rows

$$AB = \begin{bmatrix} \mathbf{a}_1 & \dots & \mathbf{a}_n \end{bmatrix} \begin{bmatrix} \mathbf{b}_1^* \\ \vdots \\ \mathbf{b}_n^* \end{bmatrix} = \mathbf{a}_1 \mathbf{b}_1^* + \dots + \mathbf{a}_n \mathbf{b}_n^*$$

Sum of rank-one matrices $\mathbf{a}_i \mathbf{b}_i^* = \mathbf{column\ times\ row}$

$$\text{Example} \quad \begin{bmatrix} 1 \\ 4 \\ 3 \end{bmatrix} \begin{bmatrix} 2 & 1 & -1 \end{bmatrix} = \begin{bmatrix} 2 & 1 & -1 \\ 8 & 4 & -4 \\ 6 & 3 & -3 \end{bmatrix}$$

Column space = all multiples of $\begin{bmatrix} 1 \\ 4 \\ 3 \end{bmatrix}$ Row space = all multiples of $\begin{bmatrix} 2 \\ 1 \\ -1 \end{bmatrix}$

Dimension of column space = 1 = Dimension of row space: Rank = 1

Basis for the column space – by example

$$A = \begin{bmatrix} 6 & 2 & 4 \\ 4 & 2 & 2 \\ 3 & 2 & 1 \end{bmatrix}$$

Column 1 is not zero – it goes in the basis

Column 2 is not a multiple of column 1 – put into the basis

Column 3 is (Column 1) – (Column 2): Dependent column

Column basis matrix $C = \begin{bmatrix} 6 & 2 \\ 4 & 2 \\ 3 & 2 \end{bmatrix}$

Column space has dimension 2 (2 vectors in the basis)

Row space has what dimension ??

Dimension of row space = Dimension of column space

Proof by factoring $A = CR$ to see **row rank = column rank**

$$\begin{array}{ccc} \begin{bmatrix} 6 & 2 & 4 \\ 4 & 2 & 2 \\ 3 & 2 & 1 \end{bmatrix} & = & \begin{bmatrix} 6 & 2 \\ 4 & 2 \\ 3 & 2 \end{bmatrix} \begin{bmatrix} 1 & 0 & 1 \\ 0 & 1 & -1 \end{bmatrix} \\ A & & C \quad R \end{array}$$

Columns of $A =$ combination of columns of C **Column basis in C**

Rows of A = combination of rows of $R =$ **Row basis in R**

RANK of A $r = 2 = 2$

$A = CMR$ has become an important factorization

Mixing matrix $M =$ invertible r by r C and R come from A

Four great factorizations

1. Symmetric $S = Q\Lambda Q^T$
eigenvectors in Q
eigenvalues in Λ
2. Every $A = U\Sigma V^T$
left singular vectors in U
singular values in Σ
right singular vectors in V
3. Orthogonalize columns $A = QR$
Orthogonal Q
Triangular R
4. Elimination $A = LU$
No row exchanges
Lower triangular L
Upper triangular U

Q has n orthonormal columns (length 1)

$$Q^T Q = \begin{bmatrix} \mathbf{q}_1^T \\ \vdots \\ \mathbf{q}_n^T \end{bmatrix} \begin{bmatrix} \mathbf{q}_1 & \cdots & \mathbf{q}_n \end{bmatrix} = \begin{bmatrix} 1 & & & \\ & 1 & & \\ & & \ddots & \\ & & & 1 \end{bmatrix} = I$$

If Q is **square** then $Q^T = Q^{-1}$ “**orthogonal matrix**” ($QQ^T = I$)

If Q is **rectangular** then QQ^T is a projection P ($QQ^T \neq I$)

$$P^2 = QQ^T QQ^T = QQ^T = P$$

Orthogonal matrices (square) are great for computation

$$\|Q\mathbf{x}\| = \|\mathbf{x}\|$$

$Q_1 Q_2$ is also orthogonal

Symmetric

$$S = Q\Lambda Q^T = \lambda_1 \mathbf{q}_1 \mathbf{q}_1^T + \cdots + \lambda_n \mathbf{q}_n \mathbf{q}_n^T$$

Positive definite if all $\lambda_i > 0$

Positive semidefinite if all $\lambda_i \geq 0$

Positive definite energy $\mathbf{x}^T S \mathbf{x} > 0$ all $\mathbf{x} \neq \mathbf{0}$

Positive semidefinite energy $\mathbf{x}^T S \mathbf{x} \geq 0$ all \mathbf{x}

Positive definite factorization $S = A^T A$ full rank A

Positive semidefinite factorization $S = A^T A$ any rank A

Key $\mathbf{x}^T S \mathbf{x} = \mathbf{x}^T A^T A \mathbf{x} = (A\mathbf{x})^T (A\mathbf{x}) \geq 0$

$$A = U\Sigma V^T = (\text{orthogonal})(\text{diagonal})(\text{orthogonal}) = \mathbf{SVD}$$

$$AV = A \begin{bmatrix} \mathbf{v}_1 & \cdots & \mathbf{v}_n \end{bmatrix} = U\Sigma = \begin{bmatrix} \mathbf{u}_1 & \cdots & \mathbf{u}_m \end{bmatrix} \begin{bmatrix} \sigma_1 & & & \\ & \ddots & & \\ & & \sigma_r & \\ & & & 0 \end{bmatrix}$$

$$\boxed{A\mathbf{v}_1 = \sigma_1\mathbf{u}_1 \quad \cdots \quad A\mathbf{v}_r = \sigma_r\mathbf{u}_r \quad A \text{ has rank } r}$$

$$A^T A = (V\Sigma^T U^T)(U\Sigma V^T) = V(\Sigma^T \Sigma)V^T$$

\mathbf{v} 's are eigenvectors of $A^T A$: orthonormal!

σ^2 are eigenvalues of $A^T A$: $\sigma^2 = \lambda \geq 0$!

\mathbf{u} 's are eigenvectors of AA^T : orthonormal!

$$\mathbf{u}_i = \frac{A\mathbf{v}_i}{\sigma_i} \quad \text{leads to} \quad \mathbf{u}_i^T \mathbf{u}_j = \frac{\mathbf{v}_i^T A^T A \mathbf{v}_j}{\sigma_i \sigma_j} = \mathbf{v}_i^T \mathbf{v}_j \frac{\sigma_j}{\sigma_i} = \begin{cases} 1 & i = j \\ 0 & i \neq j \end{cases}$$

$$A = U\Sigma V^T = \sigma_1 \mathbf{u}_1 \mathbf{v}_1^T + \cdots + \sigma_r \mathbf{u}_r \mathbf{v}_r^T \quad (\text{decreasing } \sigma_i)$$

$$A = \begin{bmatrix} \mathbf{3} & \mathbf{0} \\ \mathbf{4} & \mathbf{5} \end{bmatrix} = \frac{3}{2} \begin{bmatrix} \mathbf{1} & \mathbf{1} \\ \mathbf{3} & \mathbf{3} \end{bmatrix} + \frac{1}{2} \begin{bmatrix} \mathbf{3} & \mathbf{-3} \\ \mathbf{-1} & \mathbf{1} \end{bmatrix}$$

PCA = Principal Component Analysis uses the SVD

$$\sigma_1 \mathbf{u}_1 \mathbf{v}_1^T = \frac{3}{2} \begin{bmatrix} \mathbf{1} & \mathbf{3} \\ \mathbf{1} & \mathbf{3} \end{bmatrix} = \text{rank 1 matrix closest to } A$$

$$A_k = \sigma_1 \mathbf{u}_1 \mathbf{v}_1^T + \cdots + \sigma_k \mathbf{u}_k \mathbf{v}_k^T = \text{rank } k \text{ matrix closest to } A$$

A_k minimizes $\|A - \text{any rank } k \text{ matrix}\|$

$$\ell^2 \text{ norm } \|A\| = \max \|A\mathbf{x}\|/\|\mathbf{x}\| = \sigma_1 = \text{largest } \sigma$$

$$\text{Frobenius norm } \|A\|_F^2 = \text{sum of all } |a_{ij}|^2 = \text{sum of all } \sigma_i^2$$

$A = QR = (\text{orthogonal})(\text{triangular})$ **Gram-Schmidt**

$\mathbf{a}_1 = \mathbf{q}_1 r_{11}$ First columns of A and Q : $r_{11} = \|\mathbf{a}_1\|$

$\mathbf{a}_2 = \mathbf{q}_1 r_{12} + \mathbf{q}_2 r_{22}$ Second columns: $r_{12} = \mathbf{q}_1^T \mathbf{a}_2$ and $r_{22} = \|\mathbf{a}_2 - \mathbf{q}_1 r_{12}\|$

Every $r_{ij} = \mathbf{q}_i^T \mathbf{a}_j$ Subtract each $\mathbf{q}_i r_{ij}$ ($i < j$) from later columns \mathbf{a}_j

Version **1**: Subtract when you reach \mathbf{a}_j in step j

Version **2**: Subtract as soon as you know \mathbf{q}_i in step i

#2 allows column permutations: choose largest column in next step $i + 1$

Then columns are permuted and $AP = QR$ is numerically stable

$A = LU = (\text{lower triangular with } \ell_{ii} = 1)(\text{upper triangular})$

First row of U $\mathbf{u}_1^* = \mathbf{a}_1^* = \text{first row of } A$

First column of L $\ell_1 = (\text{first column of } A)/a_{11}$

Remove $\ell_1 \mathbf{u}_1^*$ to leave $A - \ell_1 \mathbf{u}_1^* = \begin{bmatrix} 0 & \mathbf{0}^* \\ \mathbf{0} & \mathbf{A}_1 \end{bmatrix}$ A_1 has size $n - 1$

Remove $\ell_k \mathbf{u}_k^*$ to find $A - \sum_1^k \ell_i \mathbf{u}_i^* = \begin{bmatrix} 0 & 0 \\ \mathbf{0} & \mathbf{A}_k \end{bmatrix}$ A_k has size $n - k$

Note that ℓ_k and \mathbf{u}_k^* start with $k - 1$ zeros: L and U are triangular

Finally $A = \sum_1^n \ell_i \mathbf{u}_i^* = LU = (\text{lower triangular})(\text{upper triangular})$

Ordering each of the factorizations

$$S = Q\Lambda Q^T \text{ with } \lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_n$$

$$A = U\Sigma V^T \text{ with } \sigma_1 \geq \sigma_2 \geq \dots \geq \sigma_r > 0$$

$$AP = QR \text{ with } r_{11} \geq r_{22} \geq \dots = \text{“column pivoting”}$$

$$PA = LU \text{ with all } |l_{ij}| \leq 1 = \text{“partial pivoting”}$$

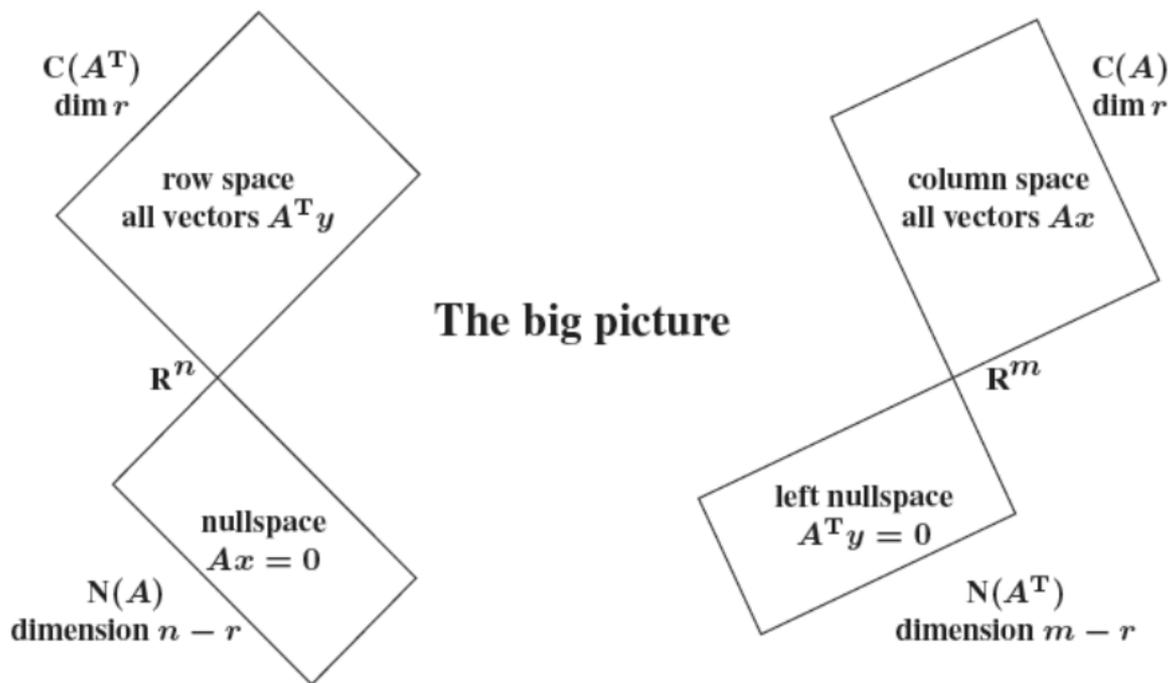
Those permutations P give numerical stability against roundo

Derivatives when $S = S(t)$ and $A = A(t)$

$$\frac{d\lambda_k(S)}{dt} = \mathbf{q}_k \frac{dS}{dt} \mathbf{q}_k^T$$

$$\frac{d\sigma_k(A)}{dt} = \mathbf{u}_k \frac{dA}{dt} \mathbf{v}_k^T$$

Four Fundamental Subspaces for A



Four Fundamental Subspaces : Their dimensions add to n and m

Pseudoinverse of $A = U\Sigma V^T$ is $A^+ = V\Sigma^+U^T$

Pseudoinverse $\Sigma = \text{diag}(\sigma_1, \dots, \sigma_r, 0, \dots, 0)$ is m by n
 $\Sigma^+ = \text{diag}(1/\sigma_1, \dots, 1/\sigma_r, 0, \dots, 0)$ is n by m

From row space to column space any A is invertible

From column space to row space A^+ is that inverse

$A^+A = \text{projection onto row space}$

$AA^+ = \text{projection onto column space}$

$A^+\mathbf{b}$ is the minimum norm least squares solution of $A\mathbf{x} = \mathbf{b}$

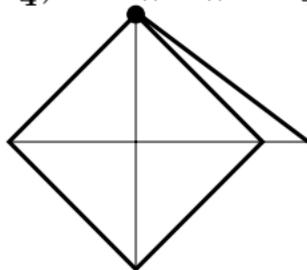
That is because $A^+\mathbf{b}$ has zero component in the nullspace of A

$A^+\mathbf{b}$ minimizes $\|A\mathbf{x} - \mathbf{b}\|^2 + \lambda\|\mathbf{x}\|^2$ as λ drops to zero

Minimization in ℓ^1 ℓ^2 ℓ^∞

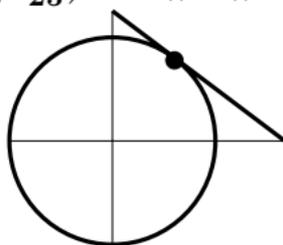
Minimize $\|v\|$ among vectors (v_1, v_2) on the line $3v_1 + 4v_2 = 1$

$(0, \frac{1}{4})$ has $\|v^*\|_1 = \frac{1}{4}$



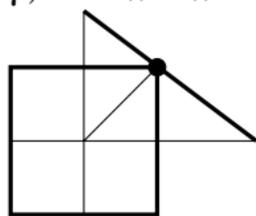
ℓ^1 diamond

$(\frac{3}{25}, \frac{4}{25})$ has $\|v^*\|_2 = \frac{1}{5}$



ℓ^2 circle

$(\frac{1}{7}, \frac{1}{7})$ has $\|v^*\|_\infty = \frac{1}{7}$



ℓ^∞ square

Basis pursuit

Minimize $|x_1| + \dots + |x_n|$ subject to $Ax = b$

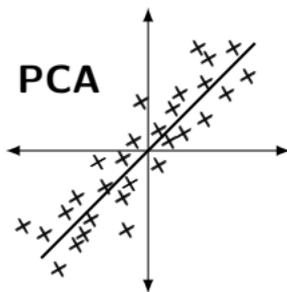
LASSO with noise

Minimize $\|Ax - b\|^2 + \lambda \sum |x_i|$

LASSO with penalty

Minimize $\|Ax - b\|^2$ with $\sum |x_i| \leq L$

Good ADMM algorithms alternate ℓ^2 problem and ℓ^1 problem



A is $2 \times n$ (large nullspace)

AA^T is 2×2 (small matrix)

$A^T A$ is $n \times n$ (large matrix)

Two singular values $\sigma_1 > \sigma_2 > 0$

The sample covariance matrix is defined by $S = \frac{AA^T}{n-1}$.

The sum of squared distances from the data points to the u_1 line is a minimum.

Total variance $T = (\sigma_1^2 + \dots + \sigma_r^2)/(n-1)$.

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